



UNIVERSITÀ
di **VERONA**

Dipartimento
di **INFORMATICA**

On the multi-interval Ulam-Rényi game: For 3 lies 4 intervals suffice

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Problem Statement

The rules of the game:

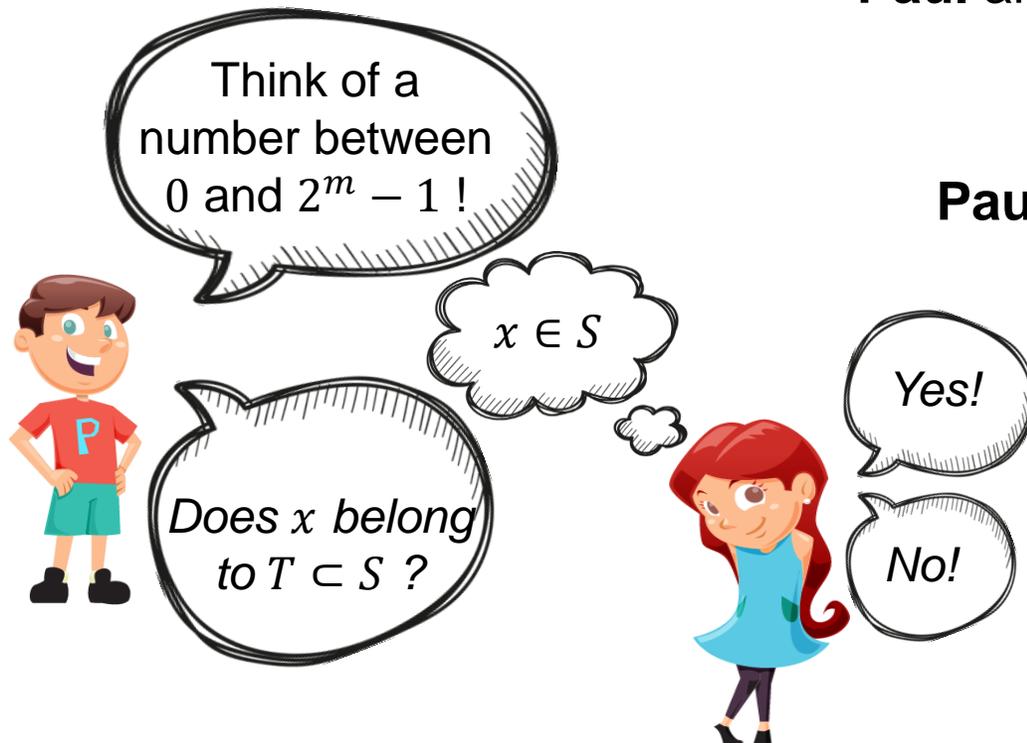
Paul and **Carole** fix a search space $S = \{0, \dots, 2^m - 1\}$
and an integer $e \geq 0$.

Carole chooses a number $x \in S$.

Paul must guess x by asking questions of the form
“does x belong to $T \subset S$?”.

Carole can answer only “yes” or “no”.

Carole can lie at most e times.



History



[1961] «... when only two things can be thought of and only one lie is allowed, then 3 questions are needed...» Rényi, A.

“On a problem of information theory”

[1968] «... at each stage the questioned set may depend on the entire past history of the game ...» Berlekamp, E. R.

“Block coding with noiseless feedback”

[1976] «... One clearly needs more than n questions for guessing one of the 2^n objects because one does not know when the lie was told...» Ulam, S. M.

“Adventures of a Mathematician”

State of the art



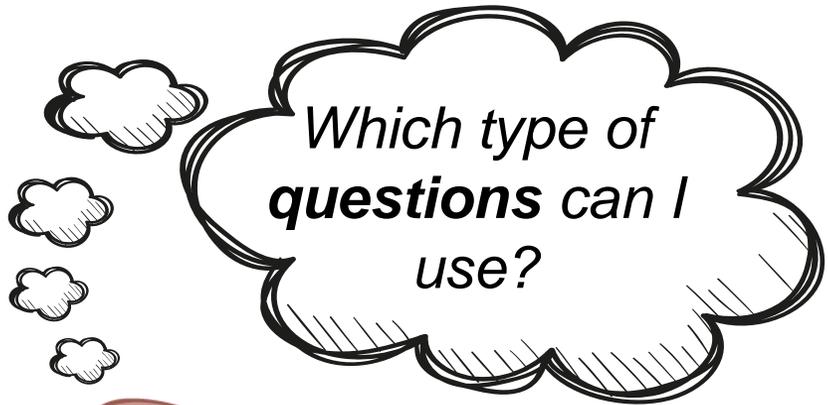
[1968] Berlekamp, E.R. Lower Bound: $\min \left\{ q \mid \sum_{j=0}^e \binom{q}{j} \leq 2^{q-m} \right\}$

Definition

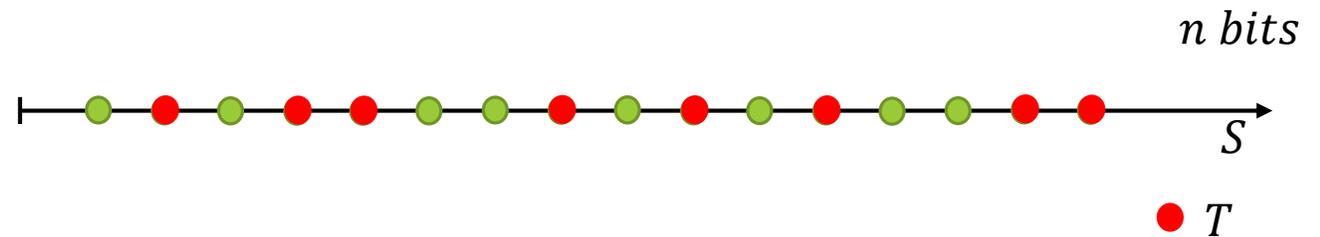
Paul's **strategy** is **perfect** if it guarantees him to identify Carole's number x using at most the number of questions given by Berlekamp's **Lower Bound**.

Number of lies (e)	Solution for $S = \{0, \dots, 2^m - 1\}$
0	Binary Search
1	[1987] Pelc, A.
2	[1989] Czyzowicz, J., Mundici, D. Pelc, A.
3	[1992] Negro, A. , Sereno, M.
e	[1992] Spencer, J.

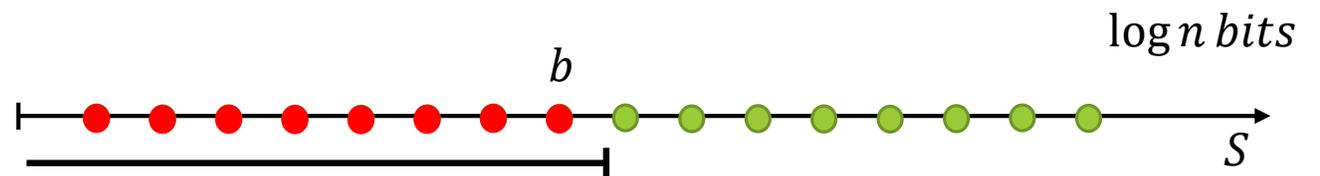
Not all subsets available



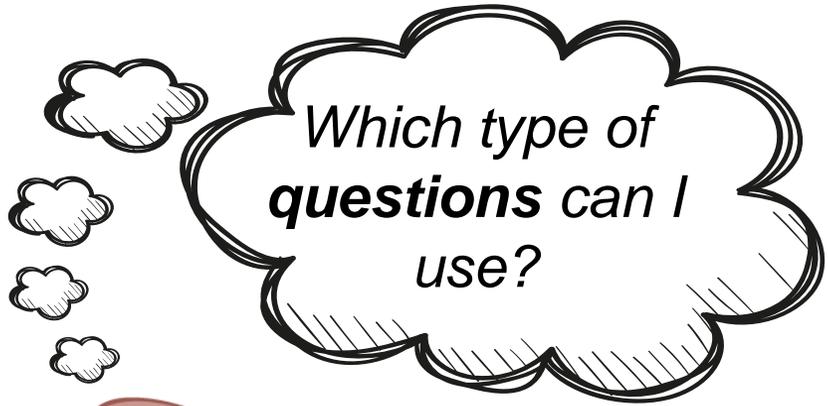
Arbitrary yes-no questions "does x belong to $T \subset S$?"



Comparison questions "Is $x \leq b$?"

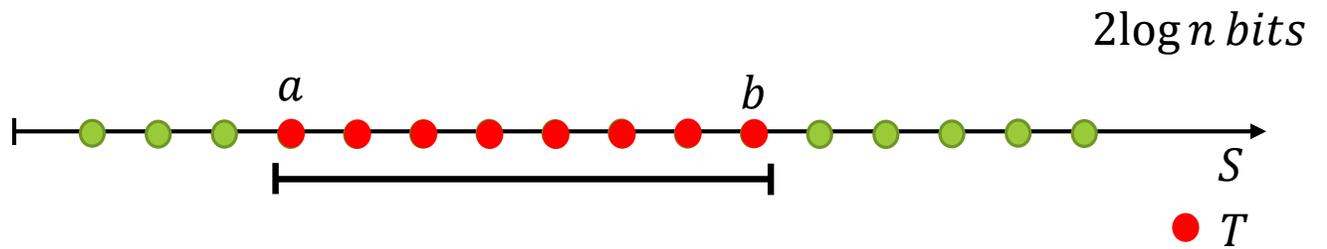


Not all subsets available

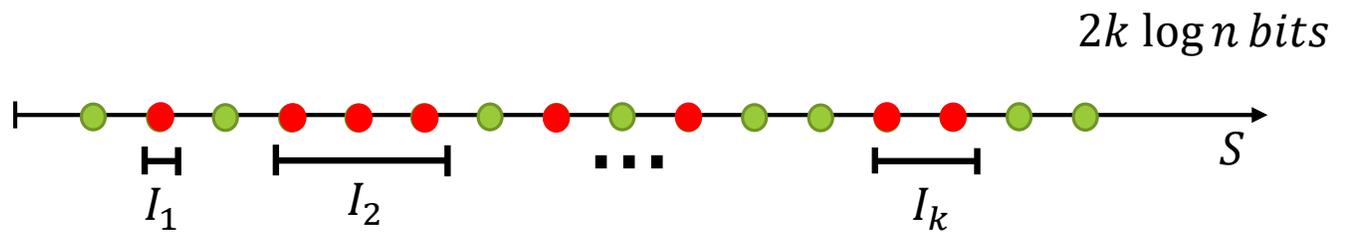


Interval questions

"Is $a \leq x \leq b$?"



Multi-Interval questions "does x belong to $I_1 \cup I_2 \cup \dots \cup I_k$?"



Main objective



Definition

Paul's **strategy** is **perfect** if it guarantees him to identify Carole's number x using at most the number of questions given by Berlekamp's **Lower Bound**.

Objective

Fixed the number of lies $e \geq 0$.

We want to find the **minimum k** such that there exists a **perfect strategy** which uses only **k-intervals questions**.

State of the art



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Number of lies (e)	Number of intervals	Solution $S = \{0, \dots, 2^m - 1\}$
0	1 / 2	<i>Binary Search</i>
1	1	[2013] Cicalese, F.
2	2	[1997] Mundici, D. Trombetta, A.
⋮	⋮	
e	$\frac{e^2+3e+2}{2}$	[2013] Cicalese, F.

Main Result



Definition

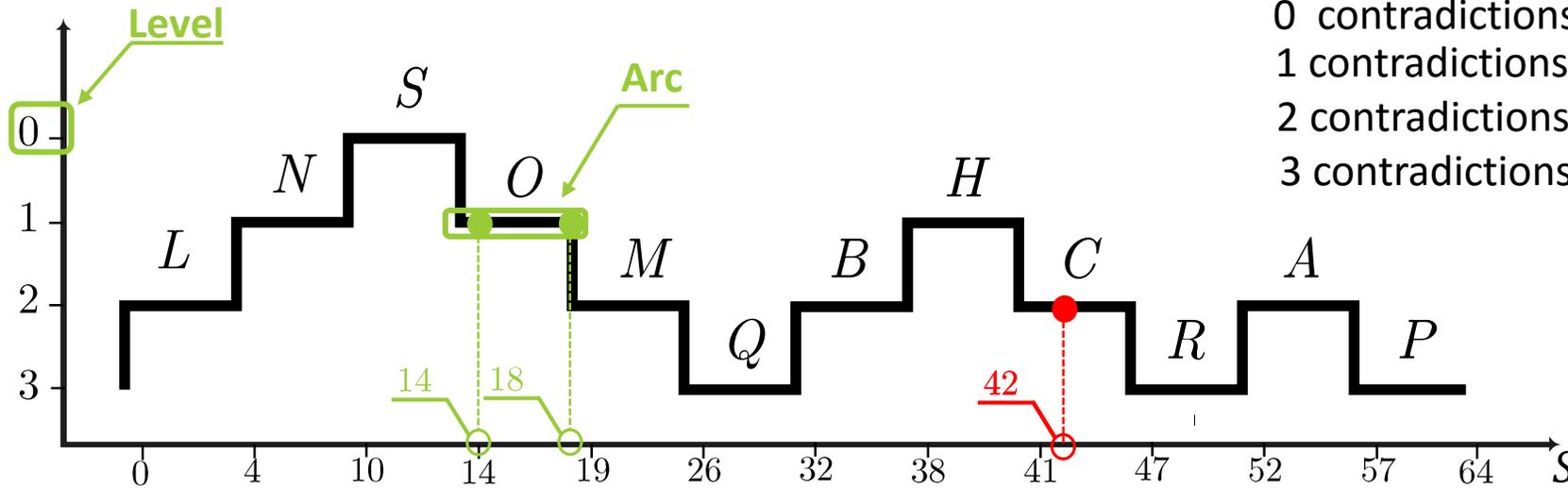
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3	4	[2017] Cicalese, F. Rossi, M.
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State Representation



- 0 contradictions: S
- 1 contradiction: $N \cup O \cup H$
- 2 contradictions: $L \cup M \cup B \cup C \cup A$
- 3 contradictions: $Q \cup R \cup P$

Structure Property:

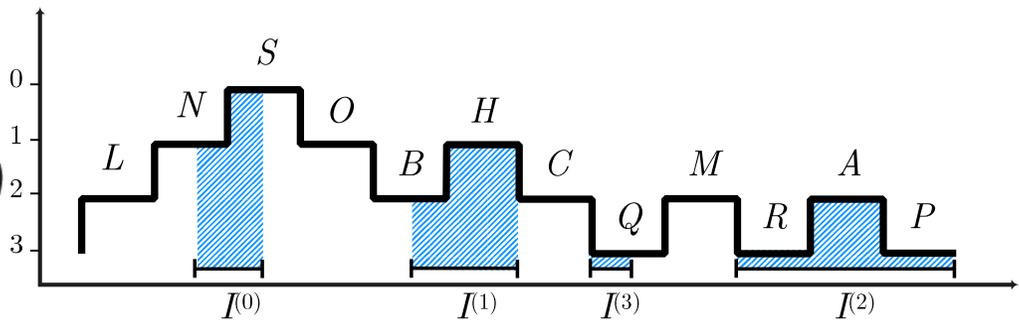
Every arc has neighboring arcs at exactly **one** level of distance.

Well-Shaped State: a state σ is well shaped iff

- For $i = 0, \dots, e - 1$, σ has exactly $2i + 1$ arcs lying on level i .
- σ has exactly e arcs lying on level e .

State Dynamics

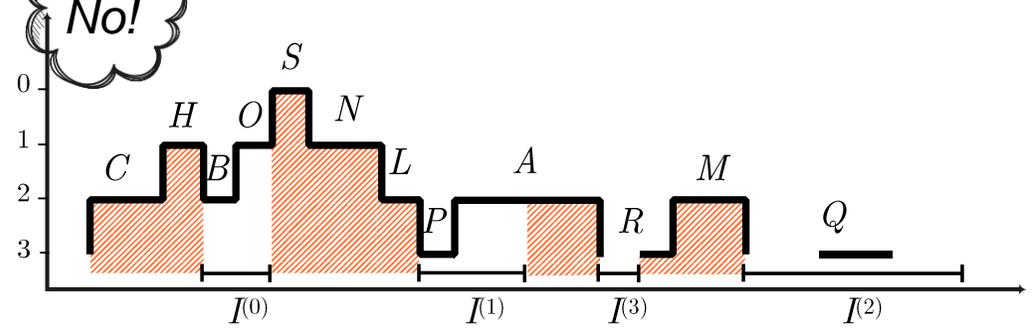
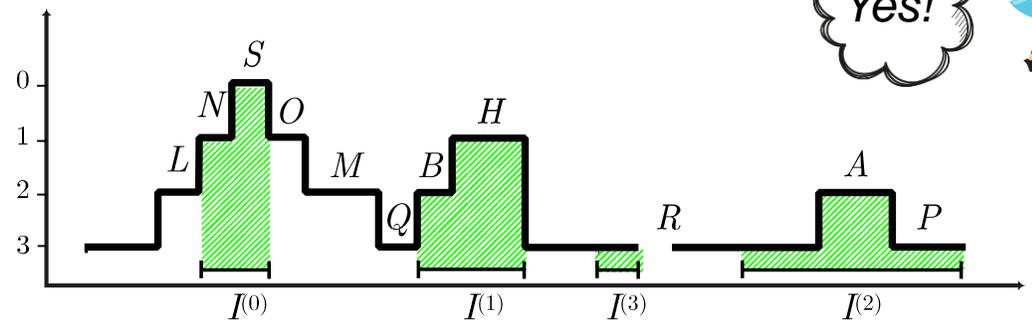
Does x belong to $I^{(0)}, I^{(1)}, I^{(2)}, I^{(3)}$?



Yes!



No!



Key Result



Definition

Given a state σ , let $ch(\sigma)$ be the minimum number of questions needed to find x according to the *Berlekamp's Lower Bound*.

Definition

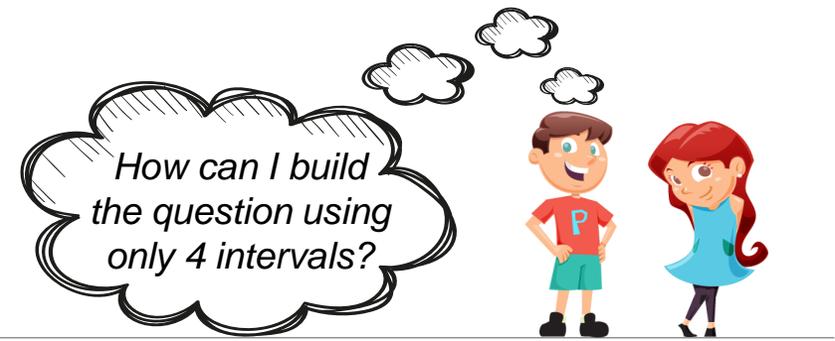
A state σ is **final** if it contains only one element, and $ch(\sigma) = 0$

Theorem

Given a well-shaped state σ there is a **4-interval question** such that the resulting states σ_{yes} and σ_{no} are well shaped and

$$\max\{ch(\sigma_{yes}), ch(\sigma_{no})\} < ch(\sigma)$$

Question Construction



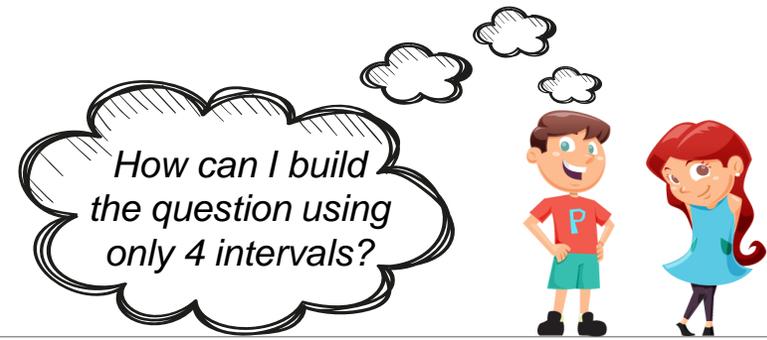
Intuition

To keep $\max\{ch(\sigma_{yes}), ch(\sigma_{no})\} < ch(\sigma)$ we **split evenly** each level between *yes* and *no*.

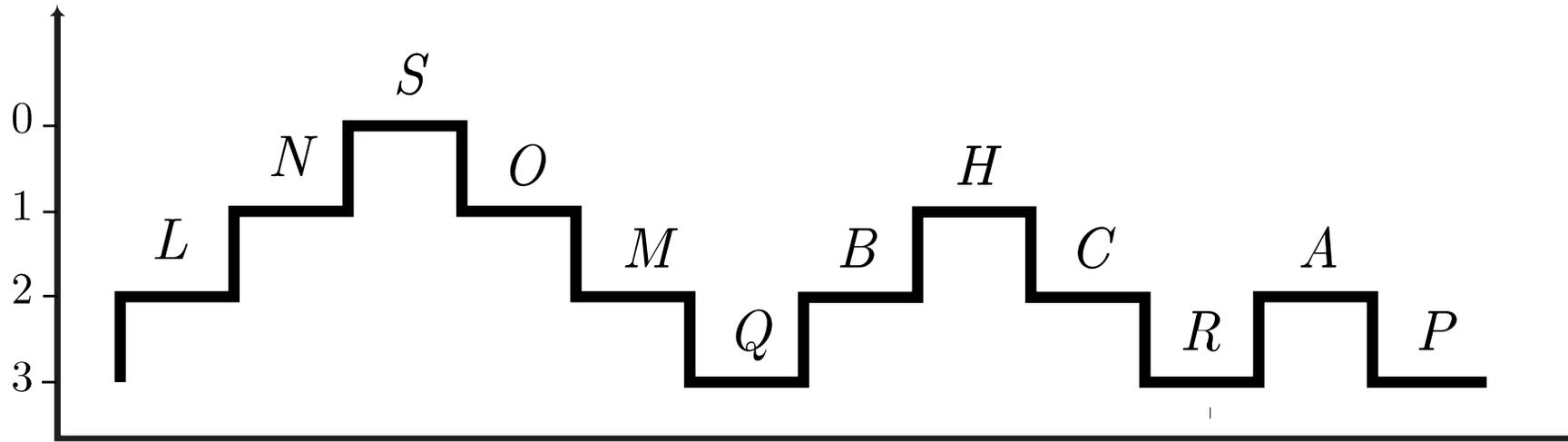
Intuition

To keep σ_{yes} and σ_{no} well-shaped we let the question **split** at most one arc per level.

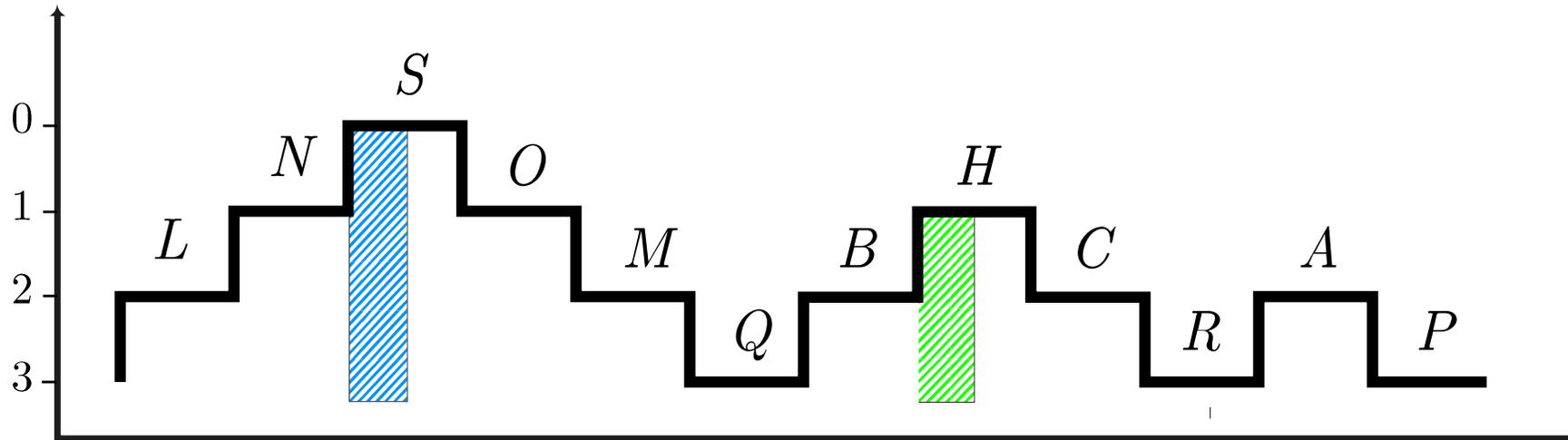
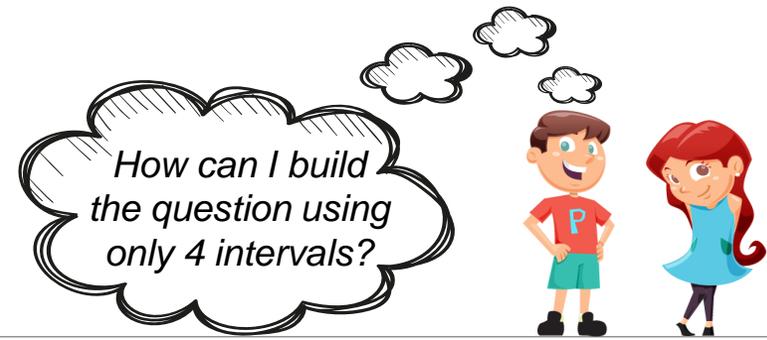
Question Construction



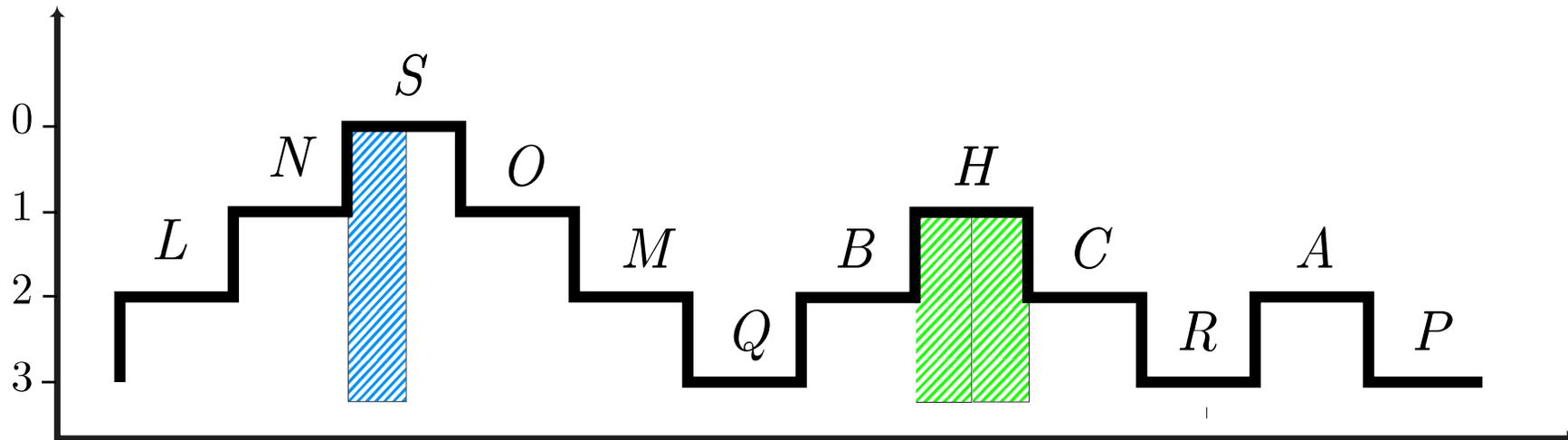
- Level 0
- Level 1
- Level 2
- Level 3



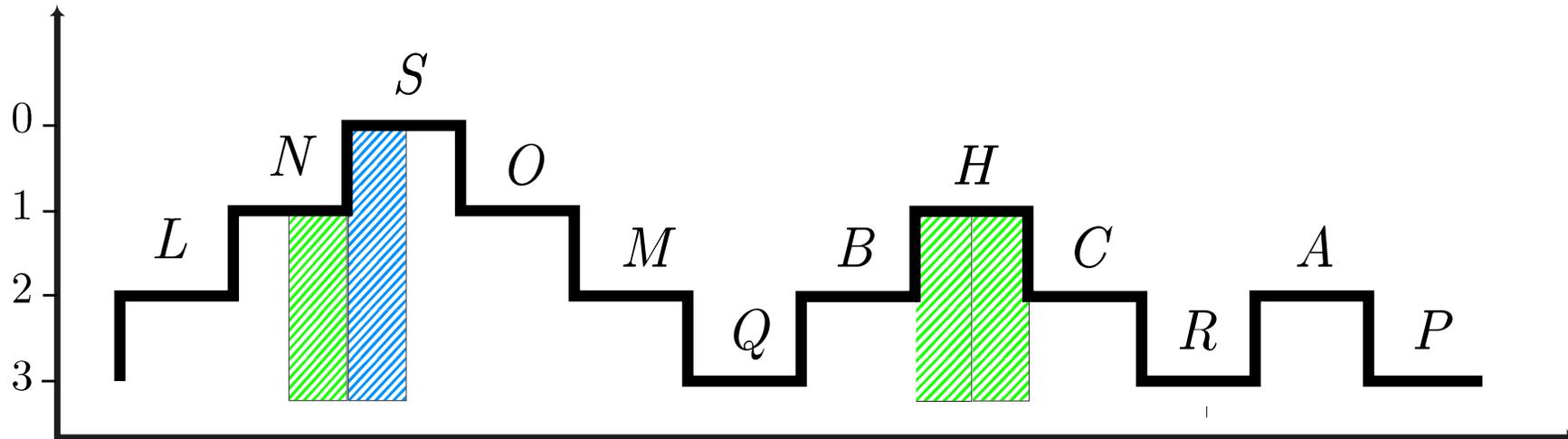
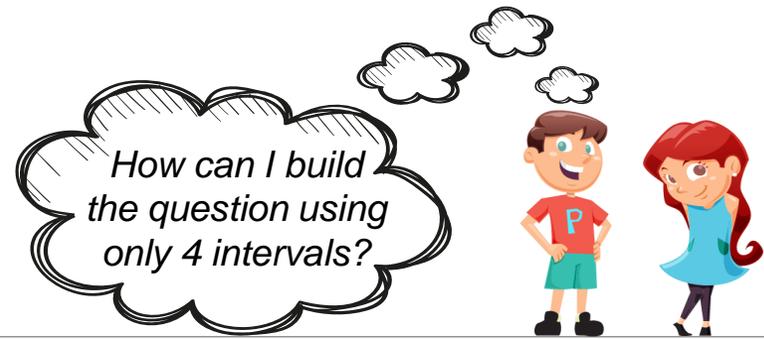
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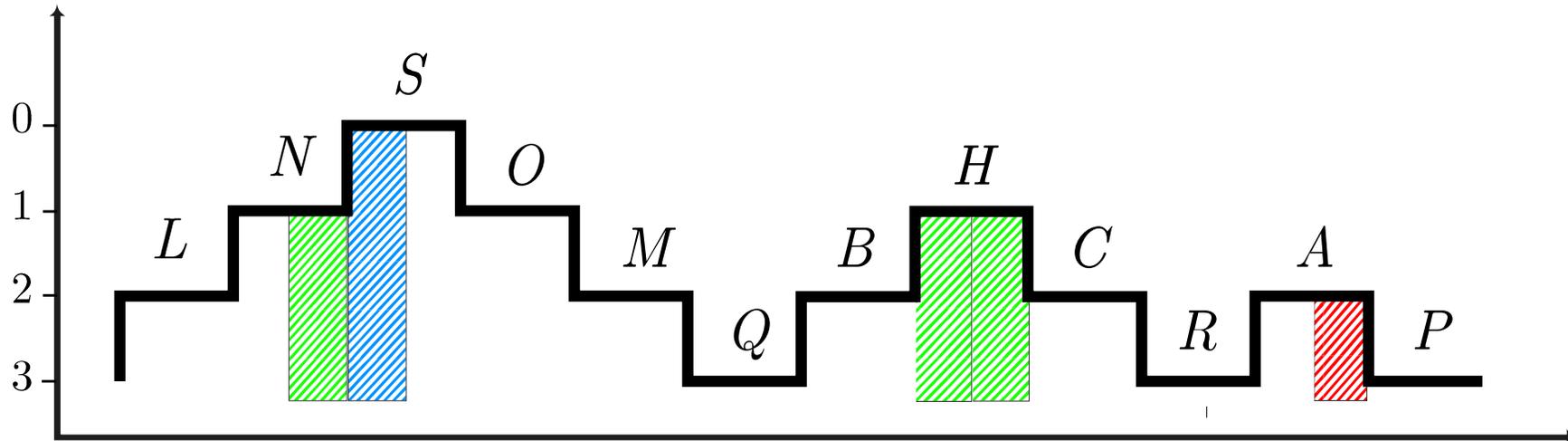
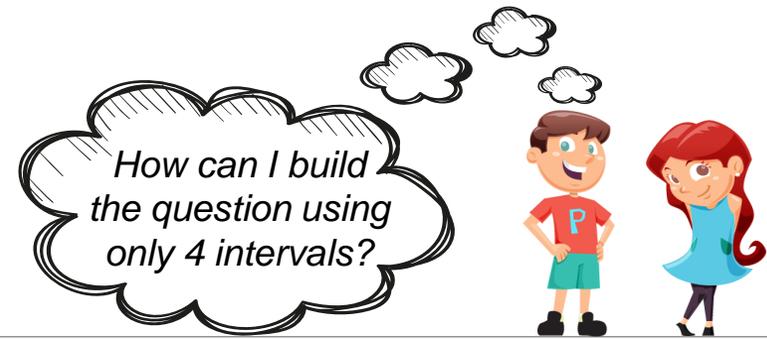
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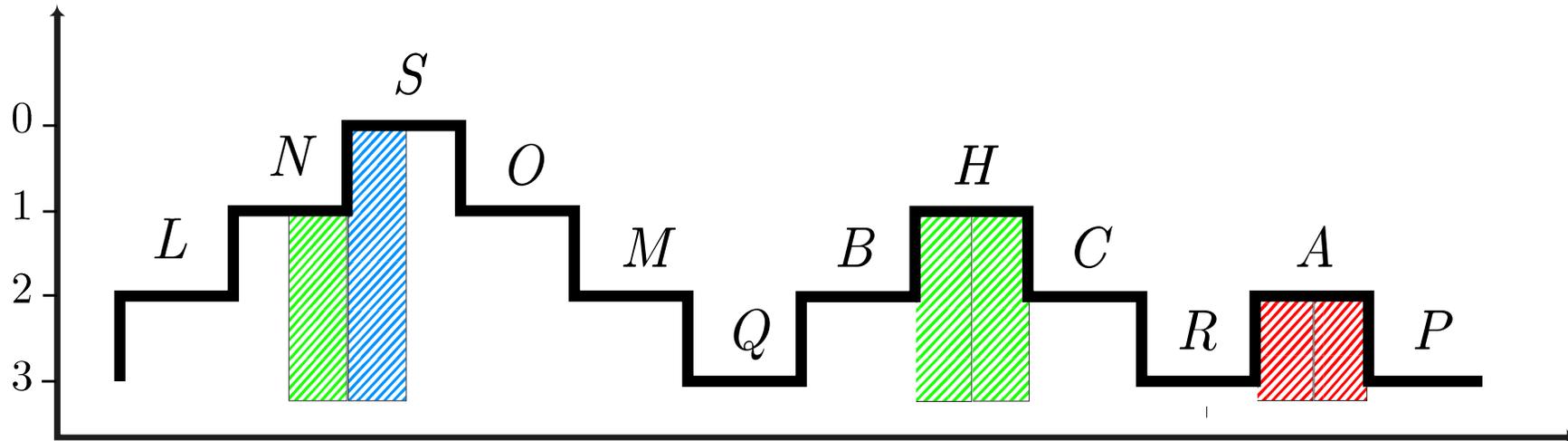
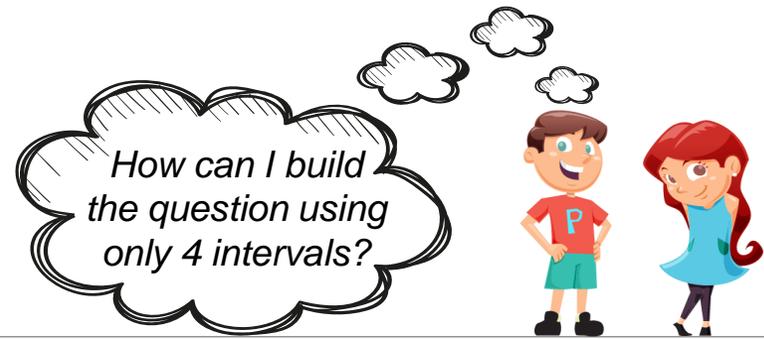
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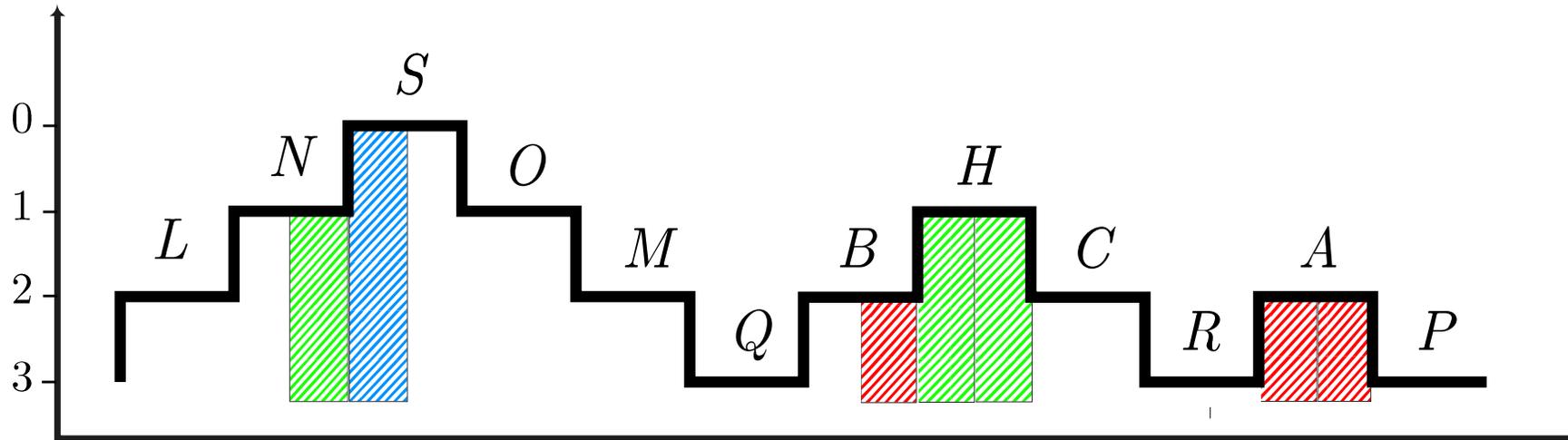
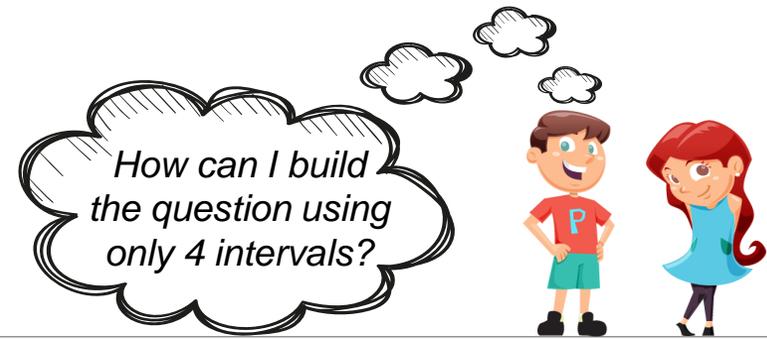
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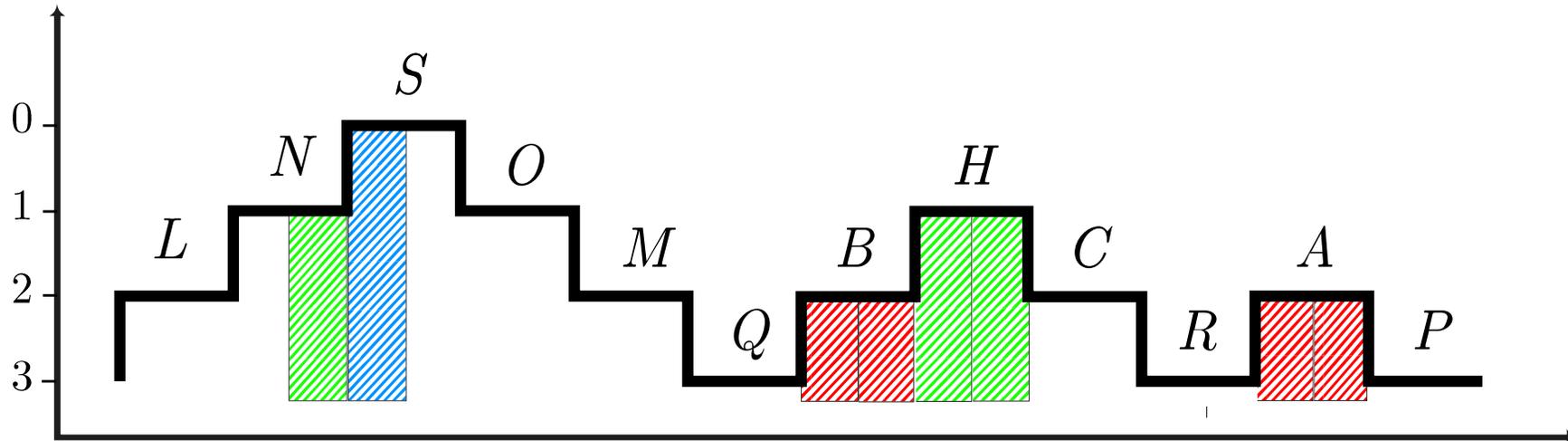
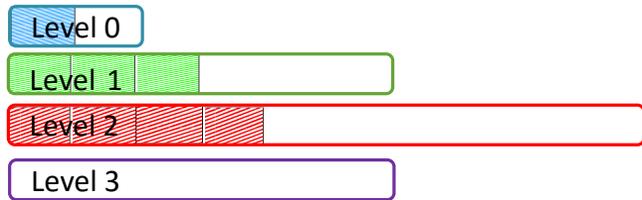
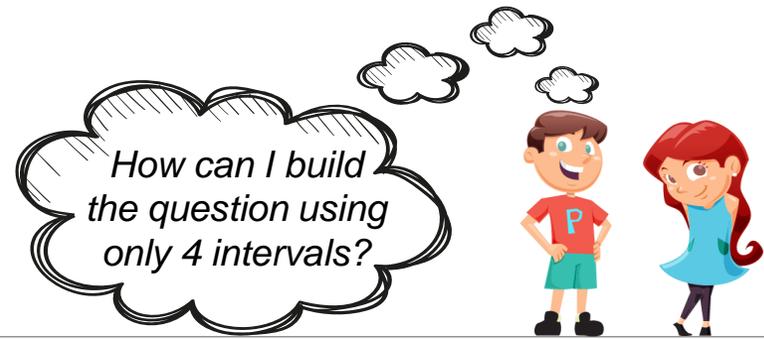
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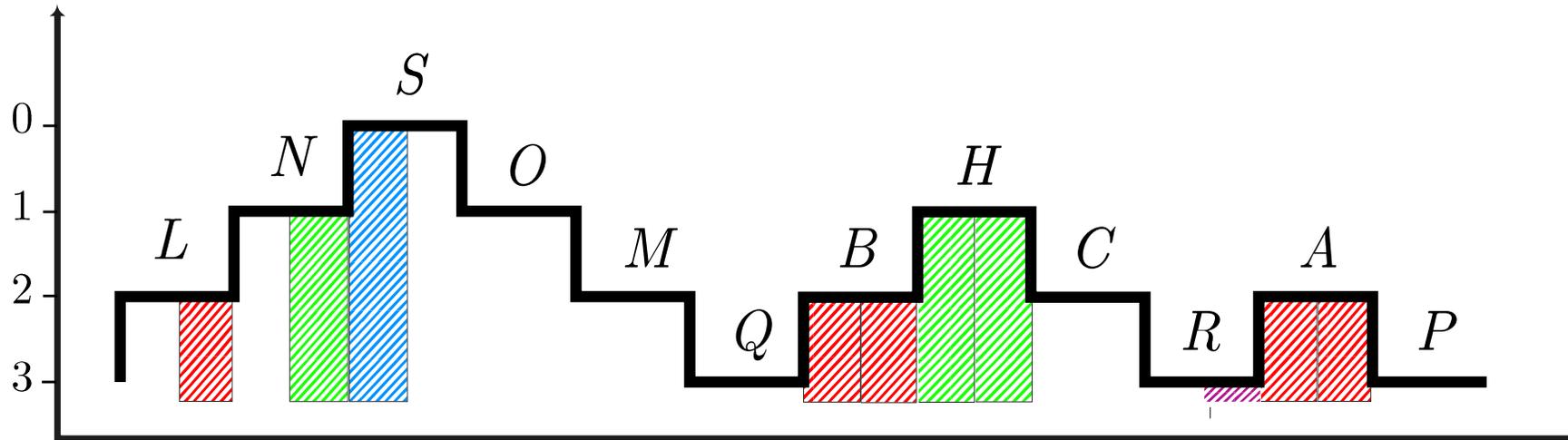
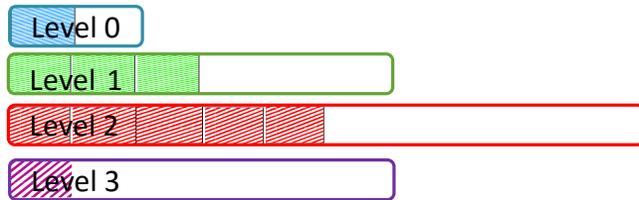
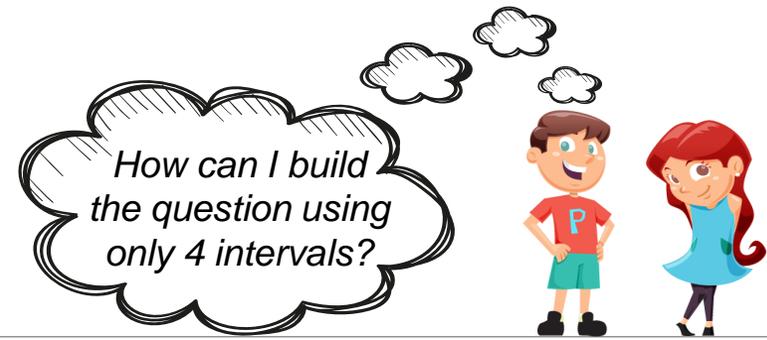
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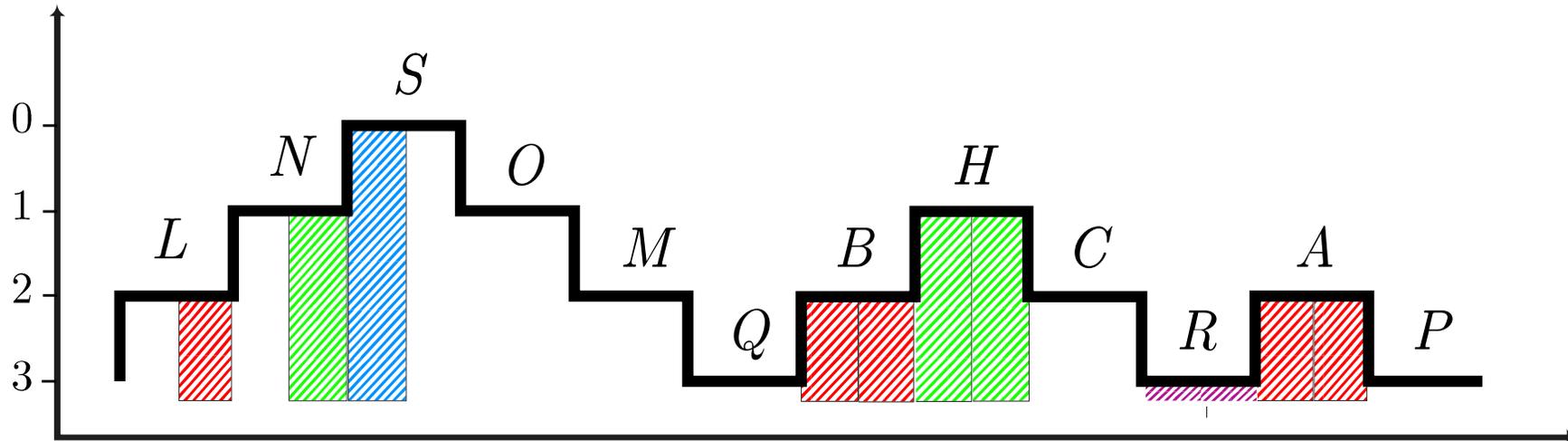
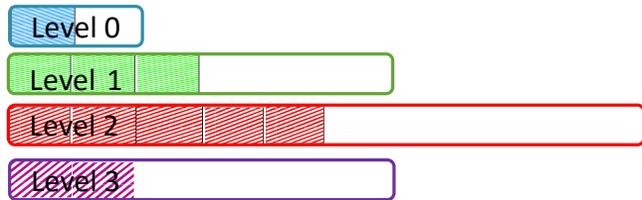
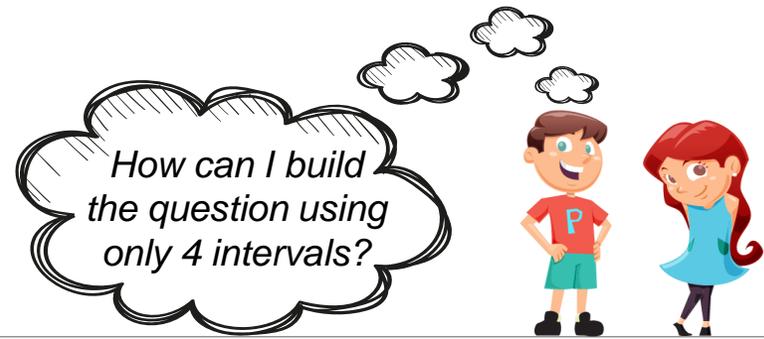
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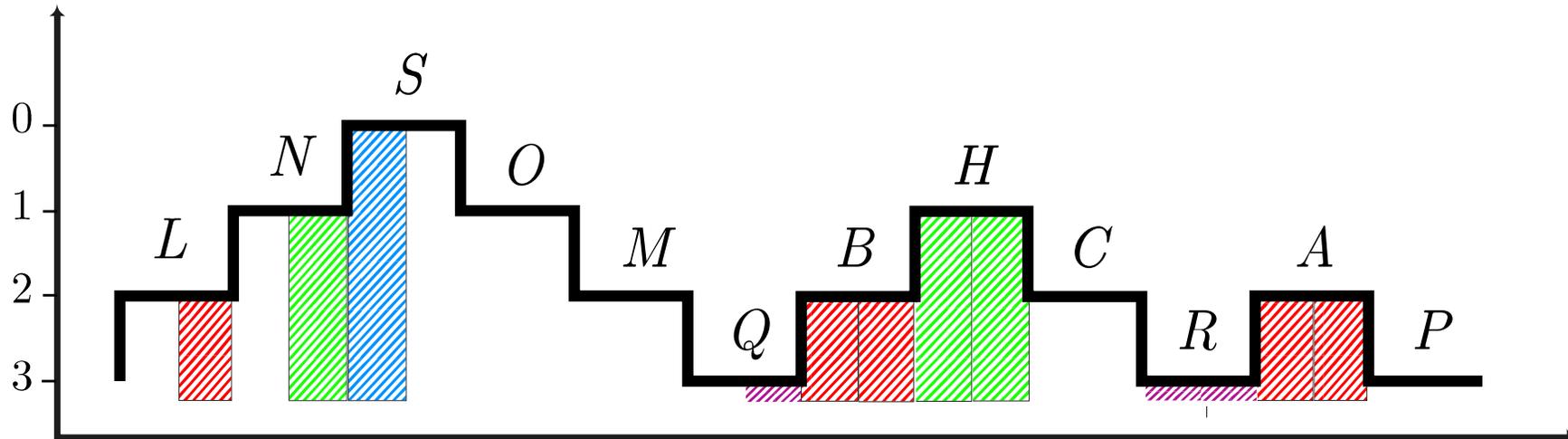
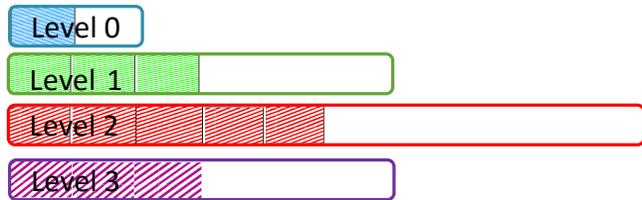
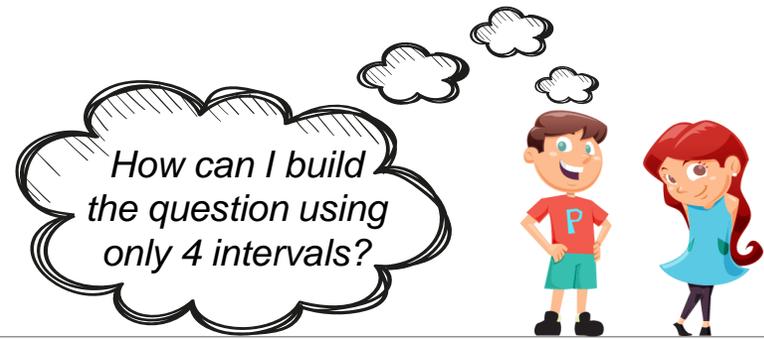
Question Construction



Question Construction



Question Construction



Conclusion



For 3 lies 4 intervals suffice.

General well-shaped structure.

Conditions for well-shapedness preservation.



Future Work



Close the gap between the bound for $e = 0,1,2,3$ and the general case $e > 3$

Non perfect strategies: more questions with less intervals per question.

Analysis of the average case.



*Thank you for
your attention!*



Questions?



Questions?

(Also non 4-intervals questions are allowed)

